

# Relational Algebra Operations

Anne Denton

Department of Computer Science  
North Dakota State University

# Outline

- 1 Basic relational algebra operations
  - Basic concepts
  - Set-theoretic operations
- 2 Cartesian product and Joins
  - Cartesian product
  - Joins
  - Recursion

# Table of Contents

- 1 Basic relational algebra operations
  - Basic concepts
  - Set-theoretic operations
- 2 Cartesian product and Joins
  - Cartesian product
  - Joins
  - Recursion

# The operations on relations produce relations

- The input and the output of all relational algebra operations are relations
- That means that arbitrary sequences of operations are possible
- For operations that are based on sets, those sets can also be the result of relational algebra operations  $\Rightarrow$  nesting
- Much of the flexibility of SQL derives from this observations

# Set theoretic operations

- Since relations are sets of tuples, any set theoretic operations can be applied
- Union, intersection, and set difference are available for relations
- Cartesian product central to capability of combining relations

# Specific operations to databases

- Select ( $\sigma$ ) selects rows
- Project ( $\pi$ ) selects columns
- Join ( $\bowtie$ ) conceptually is a combination of select, project and cartesian product

Consider the following two relations

**Student**

sid	sName
2	Black
5	Green
9	Red

**Enroll**

sid	dept	cid	grade
2	CSci	366	A
5	CSci	366	B
2	CSci	474	B

# Select

- Selects a subset of tuples from the relation
- $\sigma_{\langle \text{selection condition} \rangle}(\mathbf{R})$
- Selection condition is a Boolean expression
- Select is commutative

Result of  $\sigma_{(sid=5)}(\text{Enroll})$ :

sid	dept	cid	grade
5	CSci	366	B



## Question 1

How many rows result from the following

$\sigma_{(\text{dept}=\text{'CSci'}) \text{ AND } (\text{cid}=366)}(\text{Enroll})?$

- 1 0
- 2 1
- 3 2
- 4 3

# Project

- Projects to a subset of attributes the relation
- $\pi_{\langle \text{attribute list} \rangle}(\mathbf{R})$
- The relational algebra project eliminates duplicate rows, but the SQL implementation does not
- Project is not commutative, because attributes second project has to be a subset of those in first

Result of  $\pi_{cid}(\text{Enroll})$ :

<b>cid</b>
366
474

In a relational database all 3 entries would be preserved (366 would be listed twice)

## Brainstorming question

Why would it make sense for relational DBMSs to preserve duplicate entries?

# Table of Contents

- 1 **Basic relational algebra operations**
  - Basic concepts
  - **Set-theoretic operations**
- 2 Cartesian product and Joins
  - Cartesian product
  - Joins
  - Recursion

# Set-theoretic operations

- Relations are sets of tuples
- Anything operation can be used on sets can be used on relations
- Union:  $R \cup S$  includes tuples that are in  $R$  or in  $S$
- Intersection:  $R \cap S$  includes tuples that are in  $R$  and in  $S$
- Set Difference:  $R - S$  includes tuples that are in  $R$  but not in  $S$

Result of  $\pi_{sid}(\text{Student}) - \pi_{sid}(\text{Enroll})$ :

sid
9

## Question 2

Assume that you want to select students that take both CSci 366 and CSci 474?

Could you use  $\sigma_{(cid=474)AND(cid=366)}(Enroll)$  to produce that result?

### Question 3

Assume that you want to select students that take both CSci 366 and CSci 474?

Could you use  $\sigma_{(cid=474)}(\text{Enroll}) \cap \sigma_{(cid=366)}(\text{Enroll})$  to produce that result?

## Question 4

Assume that you want to select students that take both CSci 366 and CSci 474?

Could you use  $\pi_{sid} (\sigma_{(cid=474)}(\text{Enroll})) \cap \pi_{sid} (\sigma_{(cid=366)}(\text{Enroll}))$  to produce that result?



# Table of Contents

- 1 Basic relational algebra operations
  - Basic concepts
  - Set-theoretic operations
- 2 Cartesian product and Joins
  - Cartesian product
  - Joins
  - Recursion

# Cartesian product

- Used to combine attributes from multiple relations
- Standard set theoretic operation
- Combines each element from one with each element from another
- Usually a selection is necessary to extract the useful ones

## Question 5

How many elements does the Cartesian product of the sets  $\{1, 2, 3\}$  and  $\{A, B\}$  have

- 1 0
- 2 2
- 3 3
- 4 6

## Question 6

How many elements does the Cartesian product of the following two relations  $\{(1, A), (2, A), (3, B)\}$  and  $\{(A, x), (B, y)\}$  have?

- 1 0
- 2 2
- 3 3
- 4 6

# From Cartesian product to Join

In a database setting, one would expect the combination of the following relations

1	A
2	A
3	B

 and 

A	x
B	y

 to produce 

1	A	x
2	A	x
3	B	y

⇒ Selecting the rows with matching "join attribute" is called a natural join. After removing the redundant column the result is called natural join

# Table of Contents

- 1 Basic relational algebra operations
  - Basic concepts
  - Set-theoretic operations
- 2 Cartesian product and Joins
  - Cartesian product
  - **Joins**
  - Recursion

# Cartesian Product Example

Result of Student  $\times$  Enroll:

Student. sid	Student. sName	Enroll. sid	Enroll. dept	Enroll. cid	Enroll. grade
2	Black	2	CSci	366	A
2	Black	5	CSci	366	B
2	Black	2	CSci	474	B
5	Green	2	CSci	366	A
5	Green	5	CSci	366	B
5	Green	2	CSci	474	B
9	Red	2	CSci	366	A
9	Red	5	CSci	366	B
9	Red	2	CSci	474	B

# Making the Cartesian product useful

Superfluous rows removed by selection and columns by projection:

$$\pi_{Student.sid, Student.sName, Enroll.dept, Enroll.cid, Enroll.grade}(\sigma_{(Student.sid=Enroll.sid)}(Student \times Enroll))$$

<b>Student. sid</b>	<b>Student. sName</b>	<b>Enroll. dept</b>	<b>Enroll. cid</b>	<b>Enroll. grade</b>
2	Black	CSci	366	A
2	Black	CSci	474	B
5	Green	CSci	366	B

Result of selection and projection called a natural join:

**Student \* Enroll**



## Types of joins

- **Student**  $\bowtie$  **Enroll**: Inner join (also called equi join) does selection, but not projection and expects that join attribute specified for both
- In DBMSs the inner join is actually safer than the natural join
- **Student** \* **Enroll**: Natural join does selection and projection and assumes that join attribute has same name
- In both cases only tuples included, where join attribute appears on both sides: Student with sid=9 not included!

⇒ Motivates Outer join

# Outer Join

- **Student**  $= \bowtie$  **Enroll**: Left Outer Join includes unmatched tuples for first relation
- **Student**  $\bowtie =$  **Enroll**: Right Outer Join includes unmatched tuples for second relation
- **Student**  $= \bowtie =$  **Enroll**: Full Outer Join includes unmatched tuples for both relations

# Full Outer Join

<b>Student. sid</b>	<b>Student. sName</b>	<b>Enroll. sid</b>	<b>Enroll. dept</b>	<b>Enroll. cid</b>	<b>Enroll. grade</b>
2	Black	2	CSci	366	A
2	Black	2	CSci	474	B
5	Green	5	CSci	366	B
9	Red	<null>	<null>	<null>	<null>

## Question 7

Consider the following two relations

aid	atext
1	x
2	xx
4	y
6	yz

 and 

bid	aid
123	2
456	2
111	4
222	<null>
333	4

How many rows does the inner join have (aid is join attribute)

- 1 4
- 2 5
- 3 6
- 4 7

## Question 8

Consider the following two relations

aid	atext
1	x
2	xx
4	y
6	yz

 and 

bid	aid
123	2
456	2
111	4
222	<null>
333	4

How many rows does the left outer join have (aid is join attribute)

- 1 4
- 2 5
- 3 6
- 4 7

## Question 9

Consider the following two relations

aid	atext
1	x
2	xx
4	y
6	yz

 and 

bid	aid
123	2
456	2
111	4
222	<null>
333	4

How many rows does the right outer join have (aid is join attribute)

- 1 4
- 2 5
- 3 6
- 4 7

## Question 10

Consider the following two relations

aid	atext
1	x
2	xx
4	y
6	yz

 and 

bid	aid
123	2
456	2
111	4
222	<null>
333	4

How many rows does the full outer join have (aid is join attribute)

- 1 4
- 2 5
- 3 6
- 4 7

# Table of Contents

- 1 Basic relational algebra operations
  - Basic concepts
  - Set-theoretic operations
- 2 Cartesian product and Joins
  - Cartesian product
  - Joins
  - Recursion



## Examples where operations could be defined recursively

For example

- Employee has supervisor
- Webpage links to another

Join between a relation and itself possible but

- Every level separate join
- No recursive joining of all levels as single relational algebra command

## Question 11

Consider the following relation

aid1	aid2
1	2
1	3
2	4
4	5

Assume that you join the table with itself, joining aid2 with aid1 each time. How many rows does an inner join produce after one join?

- 0
- 1
- 2
- 4
- 6

## Question 12

Consider the following relation

aid1	aid2
1	2
1	3
2	4
4	5

Assume that you join the table with itself, joining aid2 with aid1 each time. How many rows does an inner join produce after two joins?

- 1 0
- 2 1
- 3 2
- 4 4
- 5 6

## Question 13

Consider the following relation

aid1	aid2
1	2
1	3
2	4
4	5

Assume that you join the table with itself, joining aid2 with aid1 each time. How many rows does an inner join produce after three joins?

- 1 0
- 2 1
- 3 2
- 4 4
- 5 6

## Question 14

Consider the following relation

aid1	aid2
1	2
1	3
2	4
4	5

Assume that you join the table with itself, joining aid2 with aid1 each time. How many rows does a left outer join produce after one join?

- 1 2
- 2 4
- 3 5
- 4 6

## Question 15

Consider the following relation

aid1	aid2
1	2
1	3
2	4
4	5

Assume that you join the table with itself, joining aid2 with aid1 each time. How many rows do left outer joins produce after two joins?

- 1 2
- 2 4
- 3 5
- 4 6

## Question 16

Consider the following relation

aid1	aid2
1	2
1	3
2	4
4	1

Assume that you join the table with itself, joining aid2 with aid1 each time. How many rows do left outer joins produce after one join?

- 1 2
- 2 4
- 3 5
- 4 6

## Question 17

Consider the following relation

aid1	aid2
1	2
1	3
2	4
4	1

Assume that you join the table with itself, joining aid2 with aid1 each time. How many rows do left outer joins produce after two joins?

- 1 2
- 2 4
- 3 5
- 4 6



# Problems with cycles

- Without cycles, inner joins eventually terminate
  - For inner joins no rows are left
  - For outer joins the number of rows no longer increases
- With cycles
  - The number of rows may increase indefinitely

### Question 18 (Multiple answers may be correct)

Which of the following scenarios may create indefinitely growing relations upon repeated joining?

- 1 Employees and their supervisors
- 2 Webpages linking to webpages
- 3 Publications citing earlier publications

Also: What are the respective data structures called?